

1. Motivation

- The problem of reconstruction of a matrix from an incomplete set of samples or measurements, particularly known as *matrix completion*, arises in a large area of applications including recommendation schemes, sensor network localization, collaborative filtering, quantum state tomography [1-2] etc.
- Matrix recovery based on comparisons between ratings is a very natural approach in recommendation scenarios as users are more comfortable with comparing products than giving exact ratings [3].

2. Problem Formulation

Consider a $c \times p$ rating matrix M with $[M]_{i,j} = m_{i,j}$ with rank r and c and p denoting the number of users and the number of items, respectively. We do not observe the matrix M . However, we observe a set of triplets: $\{c^{(i)}, p^{(j)}, p^{(k)}\}$ which simply illustrates a comparison: $c^{(i)}$ th user prefers $p^{(j)}$ th item over $p^{(k)}$ th item.

We then form the one-bit observation matrix $A \in \{-1, 0, 1\}^{d \times cp}$ with each of its row being a comparison.

For example:

$$M = \begin{bmatrix} 3 & 4 & 2 & 3 \\ 4 & 5 & 3 & 2 \\ 3 & 5 & 5 & 4 \end{bmatrix} \begin{matrix} c^{(1)} \\ c^{(2)} \\ c^{(3)} \end{matrix}$$

$$vec(M) = [3, 4, 3 : 4, 5, 5 : 2, 3, 5 : 3, 2, 4]^T$$

However, suppose we only have access to some comparison information:

$$\{c^{(1)}, p^{(2)}, p^{(1)}\}, \{c^{(2)}, p^{(1)}, p^{(3)}\}, \\ \{c^{(2)}, p^{(1)}, p^{(4)}\}, \{c^{(3)}, p^{(4)}, p^{(3)}\}$$

We can formulate the comparison matrix A :

$$A = \begin{bmatrix} 1 & 0 & 0 & -1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & -1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & -1 \end{bmatrix}$$

and, the one-bit comparison data λ to be:

$$\lambda = sgn(A \cdot vec(M)) = [-1 \ 1 \ 1 \ -1]^T$$

We define Ω to be the diagonalized matrix of λ ; i.e.

$$\Omega = \text{Diag}(\lambda)$$

Hence the problem of recovering the ranking matrix M will reduce to:

$$\begin{aligned} &\text{recover } M \\ \text{s. t. } &\Omega \cdot A \cdot vec(M) \geq 0 \\ &rank(M) \leq r \\ &0 \leq vec(M) \leq \eta \end{aligned}$$

Goal: To identify the low-rank matrix M , given the matrices A and Ω .

3. Proposed Approach

We expect the rating matrix to have a small rank which is a very common practice in collaborative filtering for practical reasons, we can formulate M as $M = XY^T$ and perform the alternating optimization over two tall matrices X and Y of size $c \times r$ and $p \times r$, respectively.

Thus, the matrix recovery problem can be rewritten as:

$$\begin{aligned} &\min_{M, X, Y} \|M - XY^T\|_F^2 \\ \text{s. t. } &\Omega \cdot A \cdot vec(M) \geq 0 \\ &0 < vec(M) < \eta \end{aligned}$$

which can be efficiently tackled by resorting to a cyclic minimization algorithm. The optimization problem with respect to the variable M is essentially a convex linearly-constrained quadratic program (QP), leading to a low-cost solution. Moreover, the minimizers X and Y can be obtained analytically.

$$\begin{aligned} &\|M - XY^T\|_F^2 \\ &= \|vec(M) - (Y \otimes I)vec(X)\|_F^2 \end{aligned}$$

which yields the optimal X and Y to be:

$$\begin{aligned} vec(X) &= (Y \otimes I)^\dagger vec(M) \\ vec(Y) &= (X \otimes I)^\dagger vec(M^T) \end{aligned}$$

4. The Rank-Quantization Bottleneck

$$M = x_1 y_1^T + x_2 y_2^T + \dots + x_r y_r^T$$

where $x_k \in \mathbb{R}^c$ and $y_k \in \mathbb{R}^p$. We assume that the entries of $\{x_k\}$ and $\{y_k\}$ are stored via a q -bit quantization system with a predefined set of elements and a cardinality of 2^q . As $r(c+p)q$ bits are required to store a large rating matrix in general, we need at least $r(c+p)q$ meaningful comparisons to recover M . The bottleneck is:

$$rq \leq \frac{cp - 1}{c + p}$$

5. The Rank Determination Bound

The low-rank matrix recovery algorithms will be much more effective if an initial good estimate of the matrix rank is available. Any generic row m of M is given as a linear combination of at most r vectors $\{m_k\}$.

$$m = \sum_{k=1}^r \alpha_k m_k$$

The data provide comparisons of different entries of m which can finally (or at the best performance of the system) lead to an *ordering* of the elements in m . The number of such orderings is bounded as $\mathcal{O}(n^{2r})$ which is considerably smaller than $n!$. Such a bound will help with determining a lower bound for r .

4. Results

We first consider the reconstruction of a rank-3 target rating matrix M with $c = 20$ and $p = 30$. The matrix M is generated randomly and normalized.

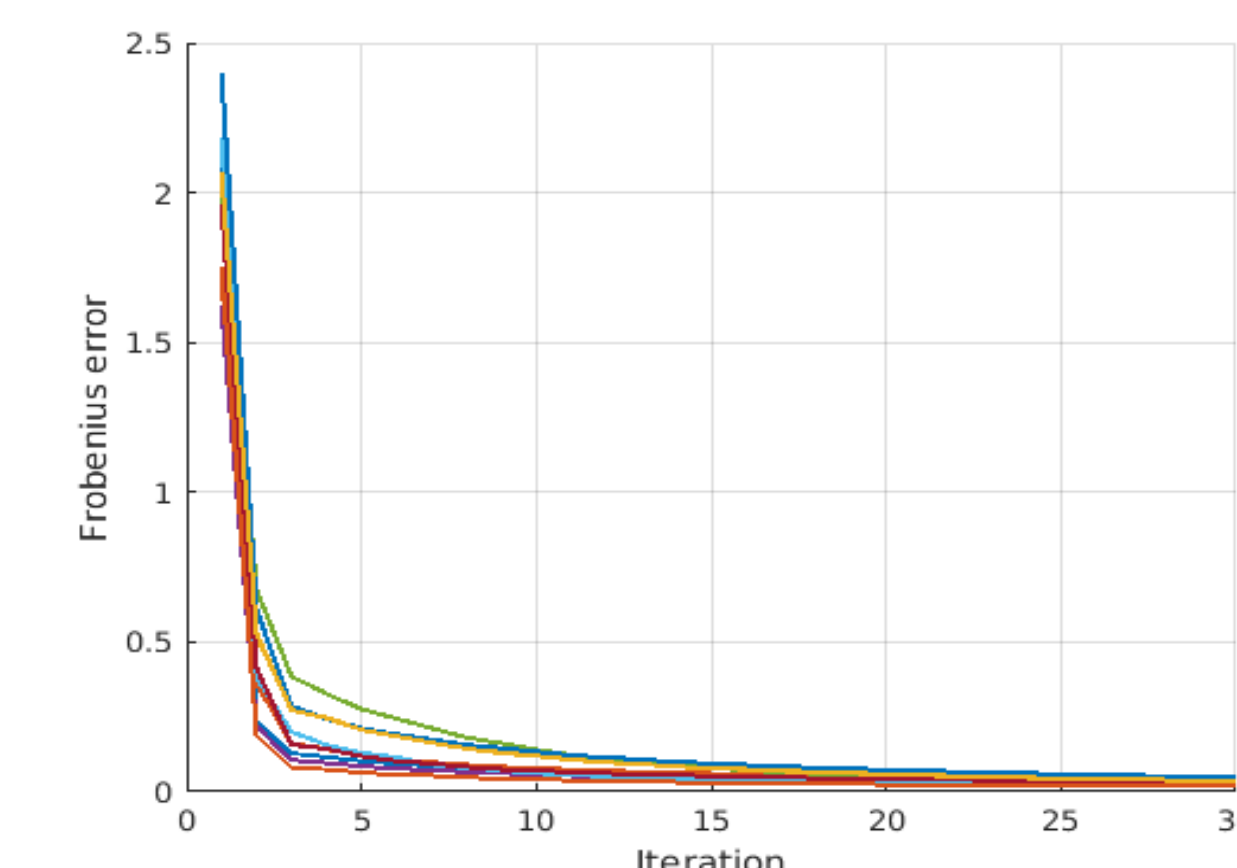


Fig. 1: Normalized Frobenius error of the low-rank matrix recovery vs. the iteration number for different random initializations with $(c; p; r) = (15; 20; 3)$

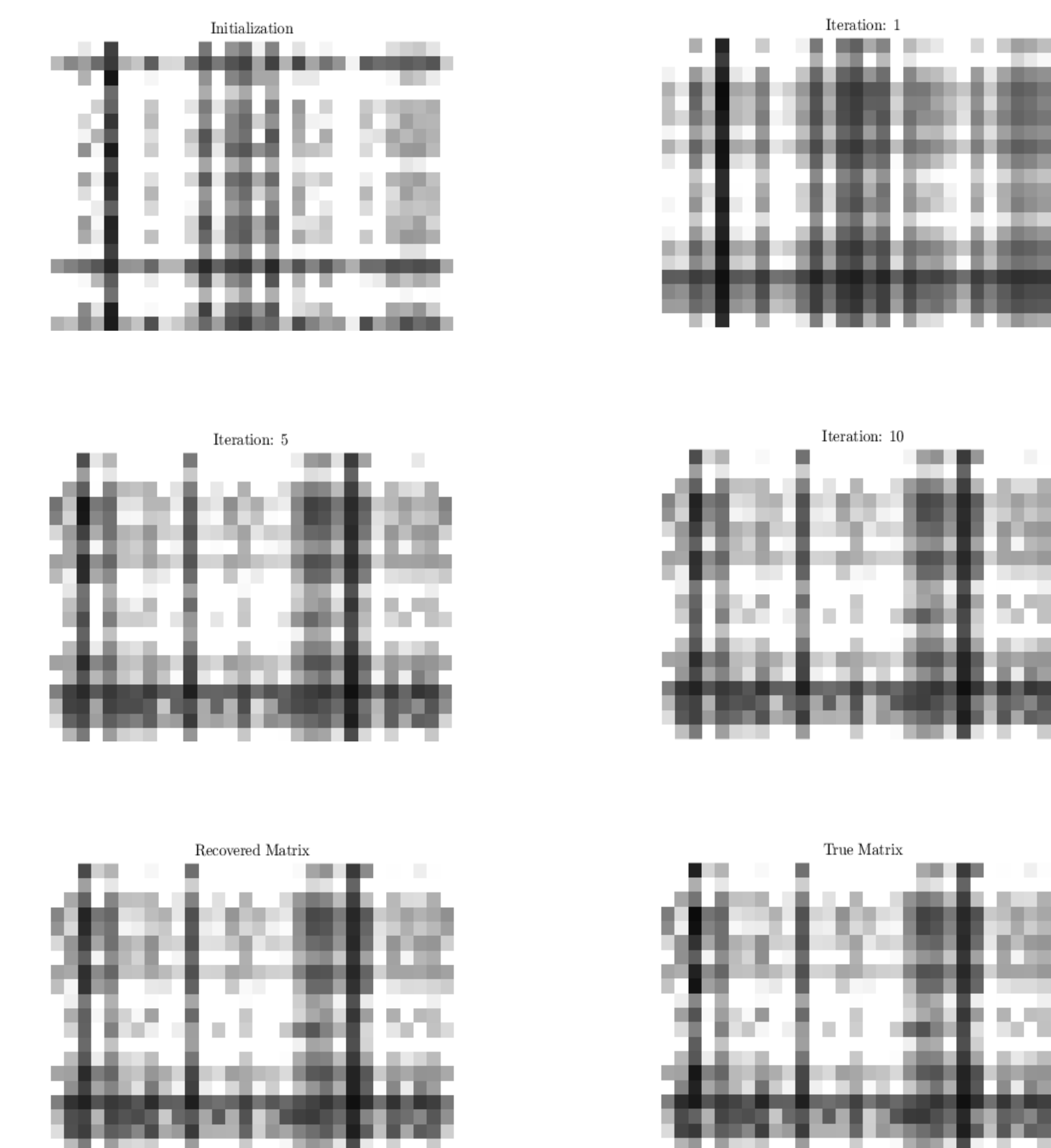


Fig. 2: An example of low-rank matrix recovery based on one-bit comparison measurements with $(c; p; r) = (15; 20; 3)$

5. References

- [1] J. F. Cai, E. J. Candes, and Z. Shen, "A Singular Value Thresholding Algorithm for Matrix Completion," *SIAM Journal on Optimization*, vol. 20, no. 4, pp. 1956–1982, Mar. 2010.
- [2] M. A. Davenport, Y. Plan, and M. V. D. Berg, E. and Wootters, "1-bit matrix completion," *Information and Inference: A Journal of the IMA*, vol. 3, no. 3, pp. 189–223, 2014.
- [3] Y. Lu, and S. N. Negahban. "Individualized rank aggregation using nuclear norm regularization." 2015 53rd Annual Allerton Conference on Communication, Control, and Computing (Allerton) (2015): 1473-1479.